



Exploiting Network Kriging for Fault Localization

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Introduction





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- Evolution toward high flexibility:
 - transmission parameters optimized for setup and changed in case of degradations/faults
 - reduction of worst-case margins to reduce costs (e.g., [a])
 - soft failures (e.g., implying Quality of Transmission QoT degradations) more frequent
 - \rightarrow need of monitoring to re-act to both soft- and hard-failures
- At the control plane, ABNO is emerging as an architecture for the control and management
 - ABNO OAM Handler responsible to receive alarms, to **correlate** the alarms, and to take actions to preserve services
- Network Kriging (NK) [b]: mathematical framework used for correlation

In this paper:

- a correlation framework based on NK for (soft or hard) fault localization
- if correlation does not solve localization (ambiguity): setup of new lightpaths with the scope of identifying unambiguously the failed elements
 - alternatively, pre-establishment of LPs for monitoring and failure localization
- We also propose a heuristic Failure Localization-Aware Routing and Spectrum Allocation (FLA-RSA) algorithm that provisions lightpaths with the objective of reducing the failure localization ambiguity



































Network Kriging concept

- Suppose LP₁, LP₂, LP₃ established and monitored
- LP₄ to be estimated
- Given additive metric y s.t. y=Gx
 - y=e2e metric, G=routing matrix, x=link metric



- $[y'_m y'_n] = [G'_m G'_n]x$, where by m we represent the lightpaths for which monitoring data are available and by n those that it should be estimated
- Can estimate \hat{y}_4 given y_1 , y_2 , y_3
- Estimation technique = "network kriging"

 $y_n = G_n G_m (G_m G^T_m)^+ y_m$





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Unambiguous:

 $AP_{AB}=1; AP_{BC}=0$





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Simulation scenario

- Compare Failure localization aware (FLA)-RSA with simple RSA
 - DT network topology
 - load is expressed as a percentage, with load=1 denoting the all-to-all communication (note that for load=1 we have unambiguous localization)
 - 100 Gbps PM-QPSK lightpaths, 37.5 GHz and 1500 km reach
- (FLA)-RSA using k=3,6,10 paths
- Metrics:
 - number of monitors as a service required for achieving unambiguous failure localization
 - number of slots required for serving traffic





Results

- Plenty extra monitors are required at low load (up to 40%) to resolve ambiguity
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- The price paid to improve failure localization is that of longer paths → a higher spectrum utilization
- The increase in spectrum is small, which is less than 5 slots for k=3





Conclusions

- We proposed a correlation framework for (soft- or hard) fault localization, leveraging information from established lightpaths
- Since a fault can be localized with ambiguity, the control plane triggers the setup of new lightpaths (monitors as a service) with the scope of identifying the failed element
- The ambiguity can be reduced using the proposed Failure Localization-Aware Routing and Spectrum Allocation (FLA-RSA) algorithm.
 - \rightarrow lower monitors as a service \rightarrow more fast and responsive reaction to failures



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